Modeling and Solving Code Generation for Real

Christian Schulte

KTH Royal Institute of Technology & SICS (Swedish Institute of Computer Science)

joint work with:	Mats Carlsson	SICS
	Roberto Castañeda Lozano	SICS + KTH
	Frej Drejhammar	SICS
	Gabriel Hjort Blindell	KTH + SICS

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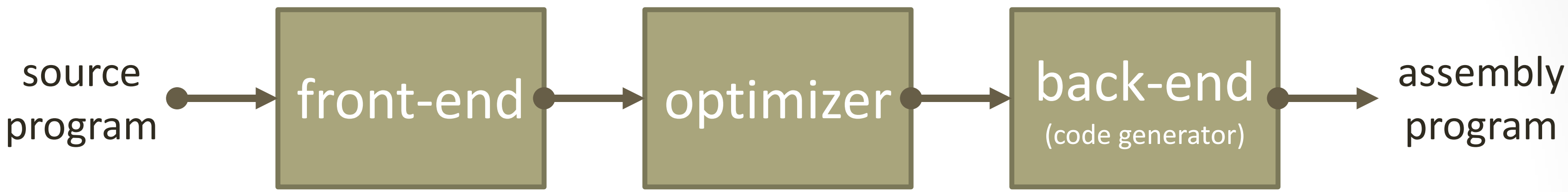
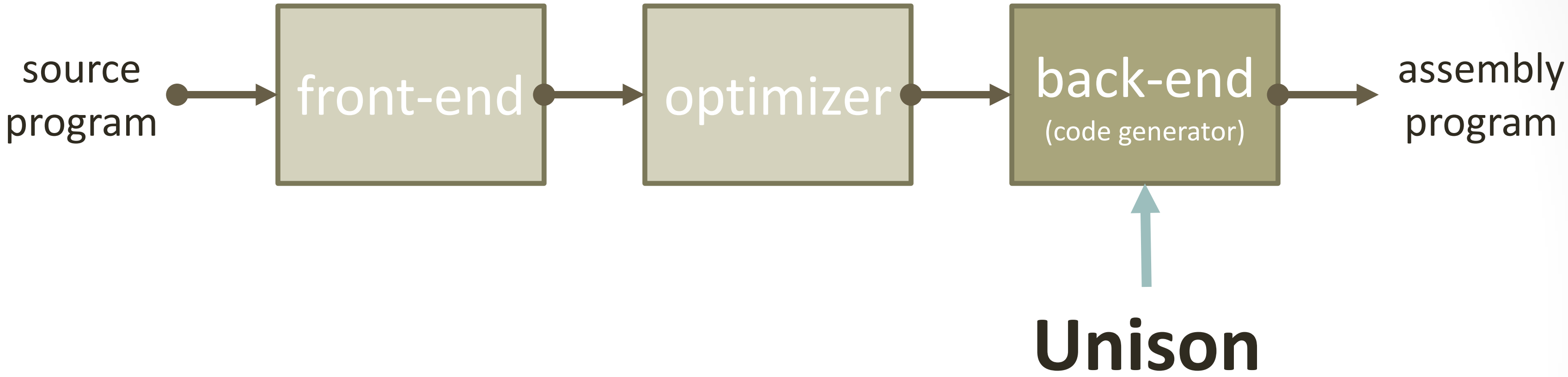
KTH Information and Communication Technology

Compilation



- Front-end: depends on source programming language
    - changes infrequently (well…)
- Optimizer: independent optimizations
    - changes infrequently (well…)
- Back-end: depends on processor architecture
    - changes often: new process, new architectures, new features, …

Generating Code: Unison



- Infrequent changes: front-end & optimizer
  - reuse state-of-the-art: LLVM, for example
- Frequent changes: back-end
  - use flexible approach: **Unison**

State of the Art

instruction selection

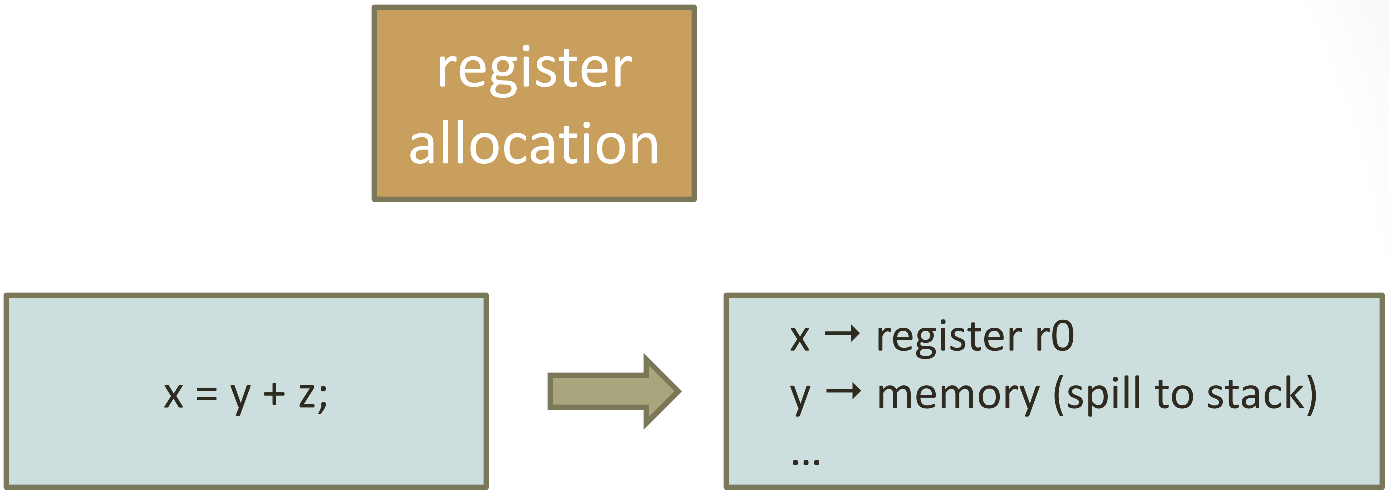
```
x = y + z;
```

→

```
add r0 r1 r2
mv  $a6f0 r0
```

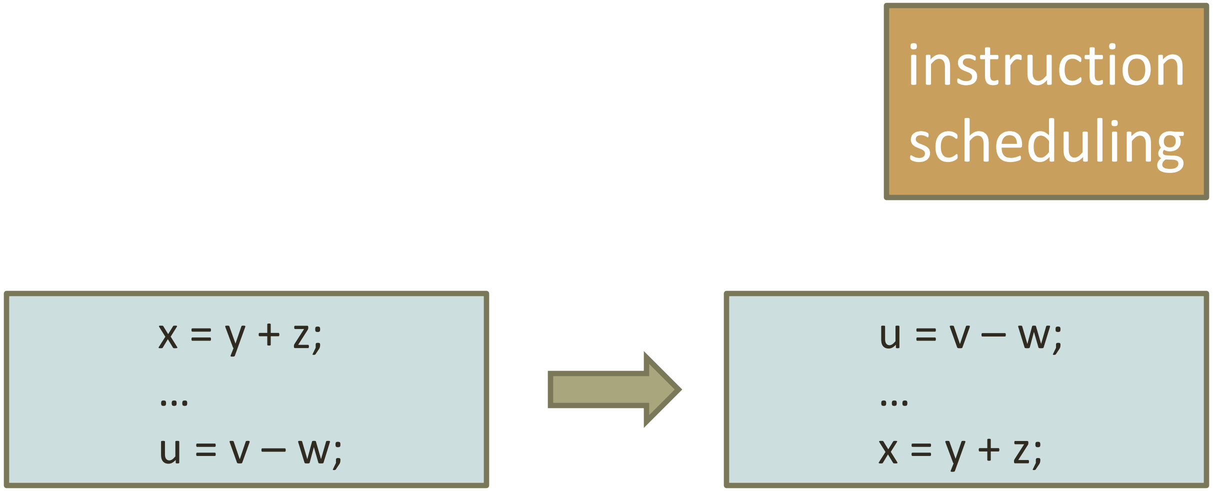
- Code generation organized into stages
    - instruction selection,

State of the Art



- Code generation organized into stages
  - instruction selection, register allocation,

State of the Art

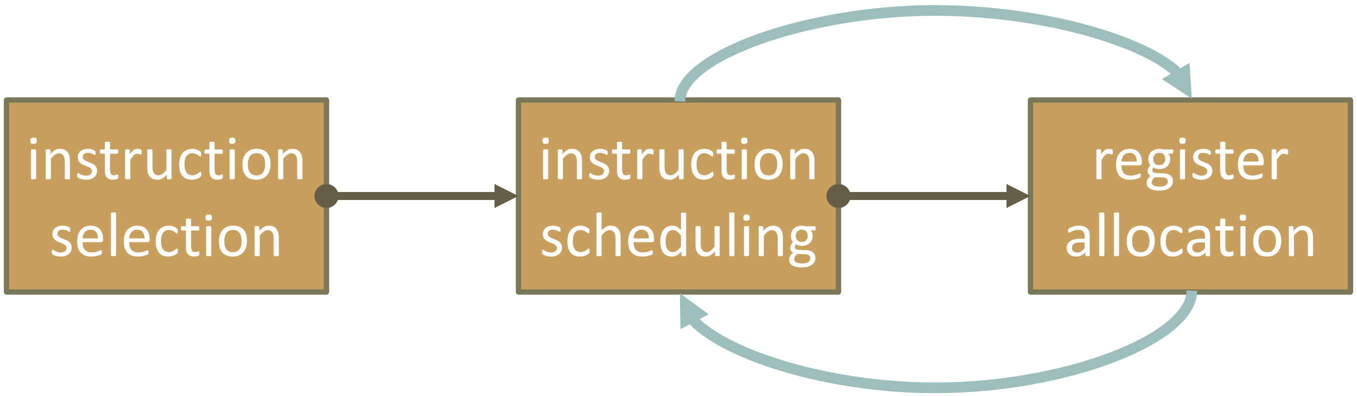


- Code generation organized into stages
  - instruction selection, register allocation, instruction scheduling

State of the Art

- Code generation organized into stages
  - stages are interdependent: no optimal order possible

State of the Art



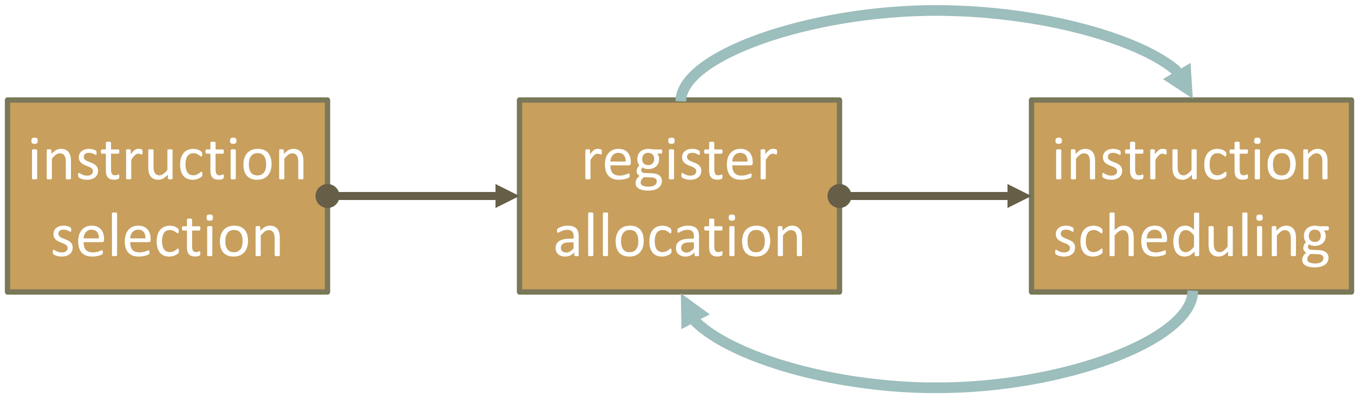
- Code generation organized into stages
  - stages are interdependent: no optimal order possible

- Example: instruction scheduling ⇆ register allocation
  - increased delay between instructions can increase throughput

    → registers used over longer time-spans

    → more registers needed

State of the Art



- Code generation organized into stages
  - stages are interdependent: no optimal order possible

- Example: instruction scheduling ⇆ register allocation
  - put variables into fewer registers

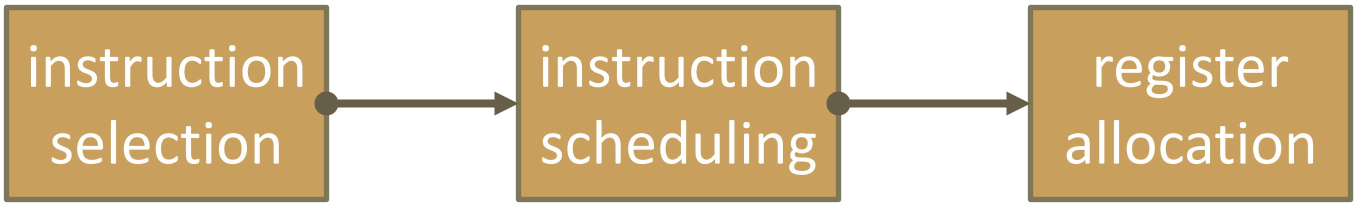
    → more dependencies among instructions

    → less opportunity for reordering instructions

State of the Art

- Code generation organized into stages
  - stages are interdependent: no optimal order possible
- Stages use heuristic algorithms
  - for hard combinatorial problems (NP hard)
  - assumption: optimal solutions not possible anyway
  - difficult to take advantage of processor features
  - error-prone when adapting to change

State of the Art



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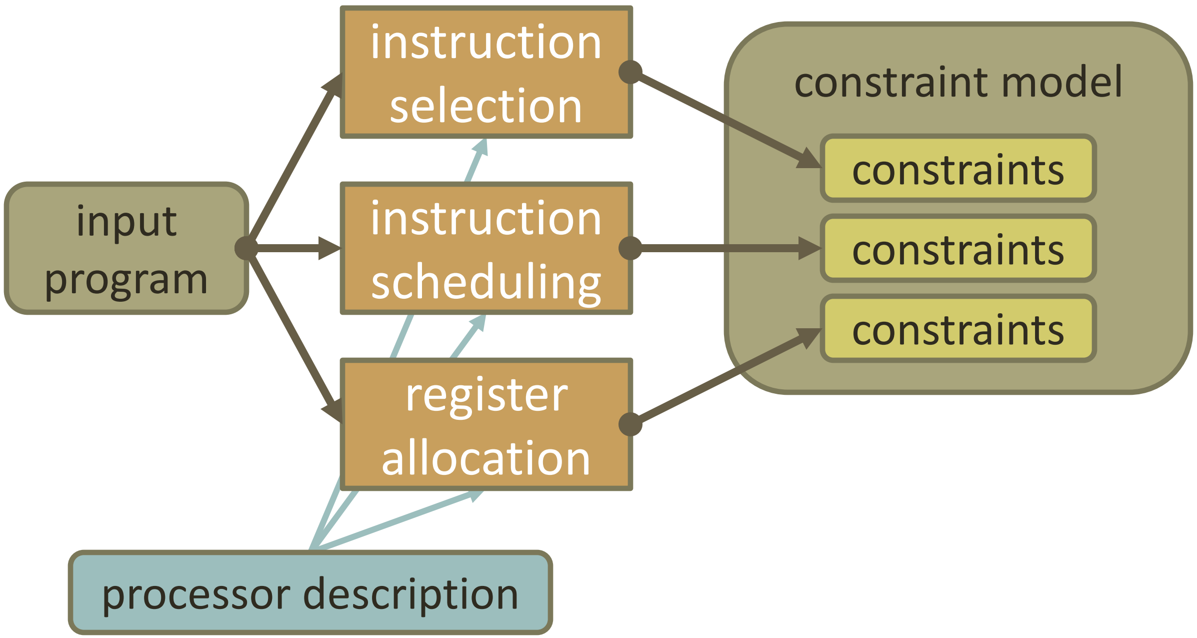
  - error-prone when adapti

preclude optimal code, make development complex

Rethinking: Unison Idea

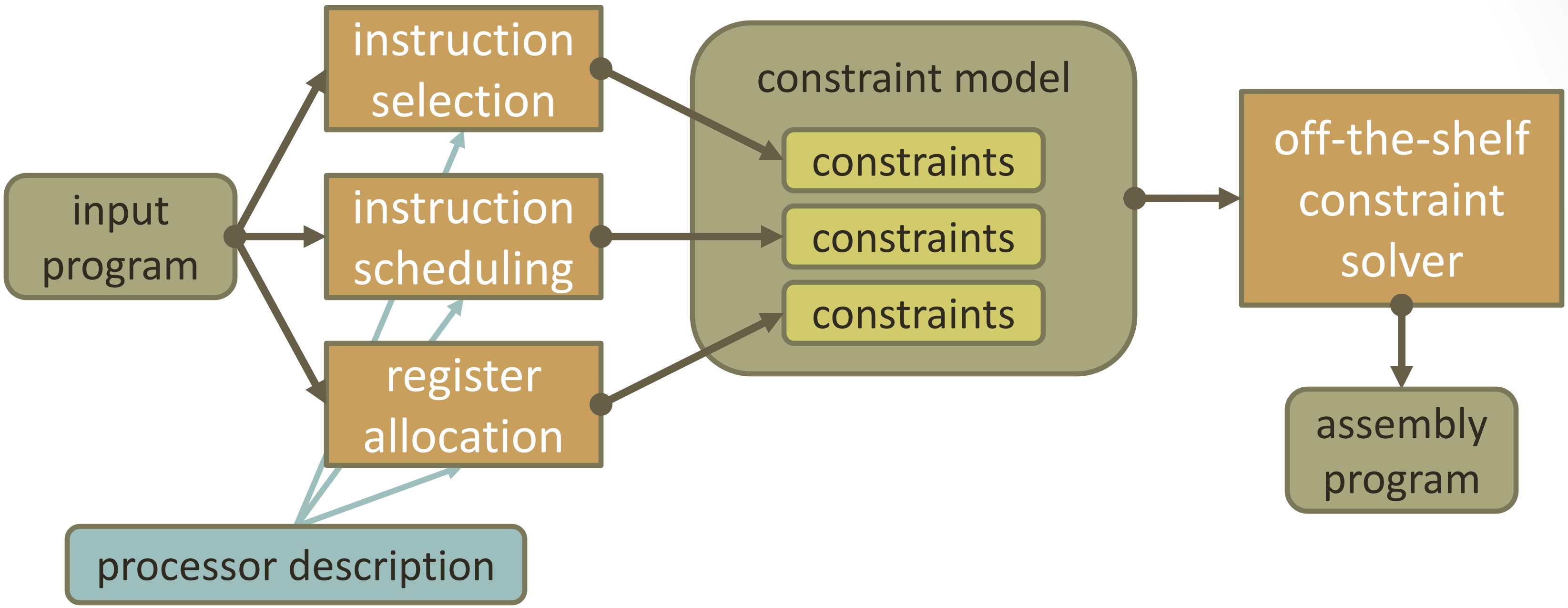
- No more staging and complex heuristic algorithms!
    - many assumptions are decades old...
- Use state-of-the-art technology for solving combinatorial optimization problems: **constraint programming**
    - tremendous progress in last two decades...
- Generate and solve single model
    - captures all code generation tasks in unison
    - high-level of abstraction: based on processor description
    - flexible: ideally, just change processor description
    - potentially optimal: tradeoff between decisions accurately reflected

Unison Approach



- Generate constraint model
  - based on input program and processor description
  - constraints for all code generation tasks
  - **generate but not solve**: simpler and more expressive

Unison Approach



- Off-the-shelf constraint solver solves constraint model
  - solution is assembly program
  - optimization takes inter-dependencies into account

Overview

- Constraint programming in a nutshell
- Register Allocation & Instruction Scheduling
  - Basic Register Allocation
  - Instruction Scheduling
  - Advanced Register Allocation [if time allows]
  - Global Register Allocation
  - Discussion
- Instruction Selection [if time allows]
  - Graph-based Instruction Selection
  - Universal Instruction Selection
  - Discussion
- Summary

CONSTRAINT PROGRAMMING IN A NUTSHELL

Constraint Programming

- Model and solve combinatorial (optimization) problems

- Modeling
  - variables
  - constraints
  - branching heuristics
  - (cost function)

- Solving
  - constraint propagation
  - heuristic search

- Of course simplified…

  …array of modeling and solving techniques

Problem: Send More Money

- Find distinct digits for letters such that

```
    SEND
+   MORE
--------
=  MONEY
```

Constraint Model

- Variables:

$$S,E,N,D,M,O,R,Y \in \{0,\ldots,9\}$$

- Constraints:

```
distinct(S,E,N,D,M,O,R,Y)

             1000×S+100×E+10×N+D
 +           1000×M+100×O+10×R+E
 = 10000×M+1000×O+100×N+10×E+Y

S≠0                M≠0
```

Constraints

- State relations between variables
  - legal combinations of values for variables

- Examples
  - all variables pair wise distinct: distinct(x_1, …, x_n)
  - arithmetic constraints: $x + 2 \times y = z$
  - domain-specific: cumulative(t_1, …, t_n)
    nooverlap(r_1, …, r_n)

- Success story: global constraints
  - modeling: capture recurring problem structures
  - solving: enable strong reasoning
    constraint-specific methods

Solving: Variables and Values

$x \in \{1,2,3,4\}$ $y \in \{1,2,3,4\}$ $z \in \{1,2,3,4\}$

- Record possible values for variables
  - solution: single value left
  - failure: no values left

Constraint Propagation

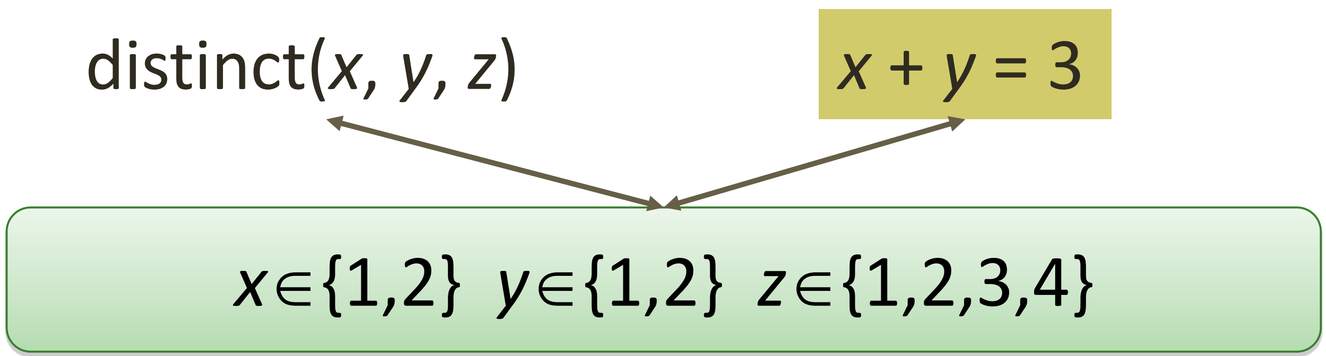
distinct(x, y, z)

$x + y = 3$

$x \in \{1,2,3,4\}$ $y \in \{1,2,3,4\}$ $z \in \{1,2,3,4\}$

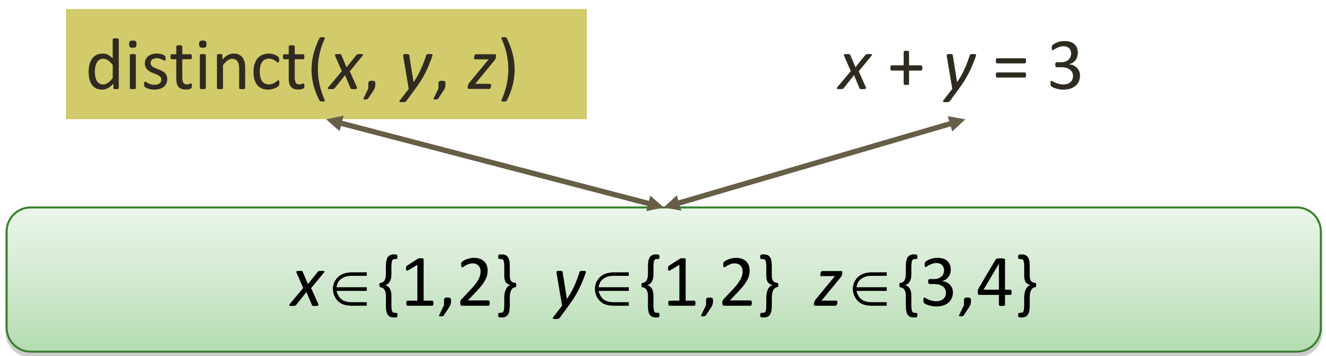
- Prune values that are in conflict with constraint

Constraint Propagation



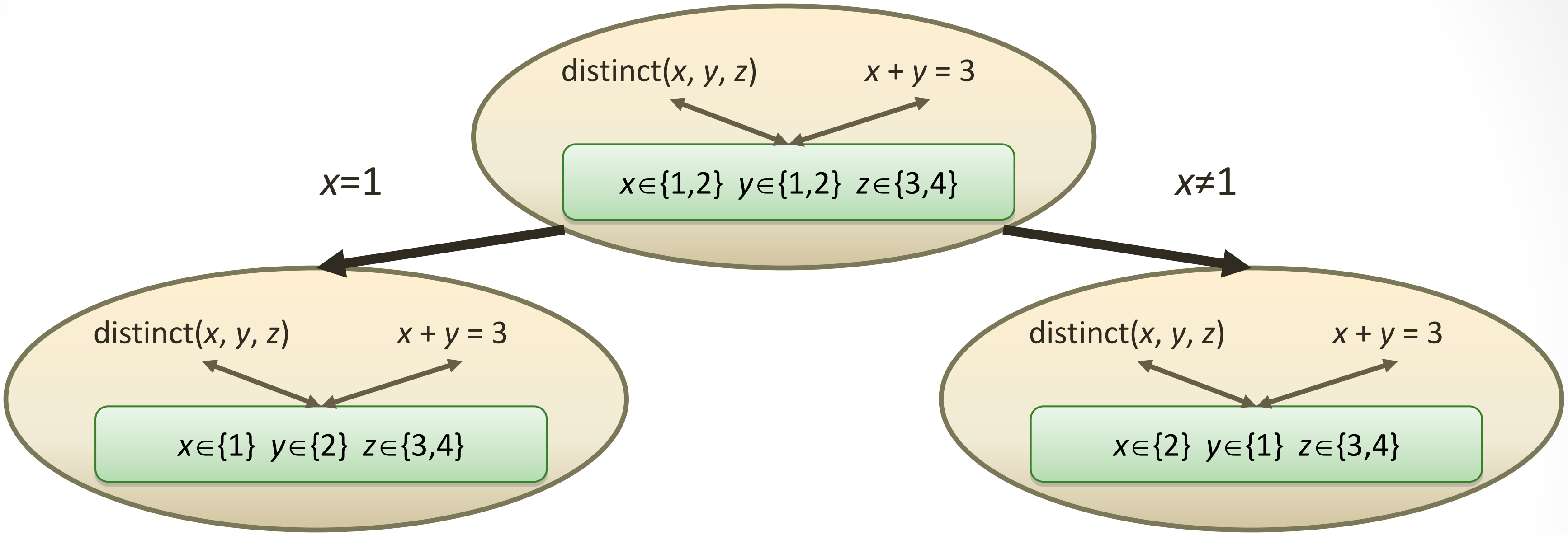
- Prune values that are in conflict with constraint

Constraint Propagation



- Prune values that are in conflict with constraint
    - propagation is often smart if not perfect!

Heuristic Search



- Propagation alone not sufficient
  - decompose into simpler sub-problems
  - search needed
- Create subproblems with additional constraints
  - enables further propagation
  - defines search tree
  - uses problem specific heuristic

What Makes It Work?

- Essential: avoid search...
  ...as it always suffers from combinatorial explosion
- Constraint propagation drastically reduces search space
- Efficient and powerful methods for propagation available
- When using search, use a clever heuristic
- Array of modeling techniques available that reduce search
- Hybrid methods (together with LP, SAT, stochastic, ...)

Register Allocation & Instruction Scheduling

Unit and Scope

- Function is unit of compilation
  - generate code for one function at a time
- Scope
  - **local** generate code for each basic block in isolation
  - **global** generate code for whole function
- **Basic block**: instructions that are always executed together
  - execute at start
  - execute all instructions
  - leave execution at end
  - that is: no control flow within basic block (in or out)

Local (and slightly naïve) register allocation

BASIC REGISTER ALLOCATION

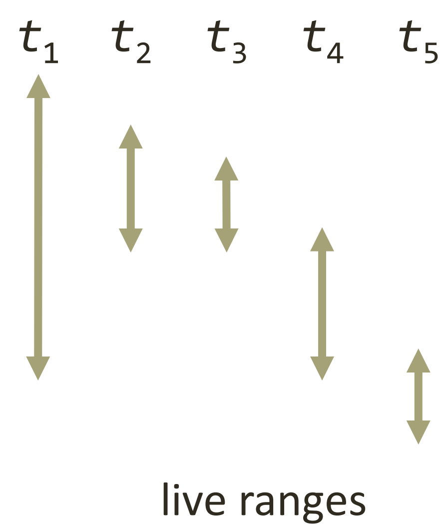
Local Register Allocation

```
t₂ ← mul t₁, 2
t₃ ← sub t₁, 2
t₄ ← add t₂, t₃
     ...
t₅ ← mul t₁, t₄
   ← jr t₅
```

- Instruction selection has already been performed
- Temporaries
  - **defined** or **def**-occurence (lhs) t_3 in $t_3 \leftarrow$ sub t_1, 2
  - **used** or **use**-occurence (rhs) t_1 in $t_3 \leftarrow$ sub t_1, 2
- Basic blocks are in SSA (single static assignment) form
  - each temporary is defined once
  - standard state-of-the-art approach

Liveness & Interference

```
t₂ ← mul t₁, 2
t₃ ← sub t₁, 2
t₄ ← add t₂, t₃
      ...
t₅ ← mul t₁, t₄
   ← jr t₅
```



- Temporary is **live** from def to last use, defining its **live range**
  - live ranges are **linear** (basic block + SSA)
- Temporaries **interfere** if their live ranges overlap
- Non-interfering temporaries can be assigned to same register

Spilling

- If not enough registers available: **spill**
- Spilling moves temporary to memory (stack)
  - store in memory after defined
  - load from memory before used
  - memory access typically considerably more expensive
  - decision on spilling crucial for performance
- Architectures might have more than one register bank
  - some instructions only capable of addressing a particular bank
  - “spilling” from one register bank to another
- **Unified register array**
  - limited number of registers for each register file
  - memory is just another “register” file
  - unlimited number of memory “registers”

Coalescing

- Temporaries *d* ("destination") and *s* ("source") are **move-related** if

  $d \leftarrow s$

  - *d* and *s* should be **coalesced** (assigned to same register)
  - coalescing saves move instructions and registers

- Coalescing is important due to
  - how registers are managed (calling convention)
  - how our model deals with global register allocation (more later)

Copy Operations

- Copy operations replicate a temporary t to a temporary t'

  $t' \leftarrow \{i_1, i_2, \ldots, i_n\}\ t$

  - copy is implemented by one of the alternative instructions $i_1, i_2, \ldots, i_n$
  - instruction depends on where t and t' are stored

    similar to [Appel & George, 2001]

- Example MIPS32

  $t' \leftarrow$ {`move`, `sw`, `nop`} t

  - t' memory and t register: `sw` spill
  - t' register and t register: `move` move-related
  - t' and t same register: `nop` coalescing
  - MIPS32: instructions can only be performed on registers

Model Variables

- Decision variables
  - reg(t) $\in$ **N** register to which temporary t is assigned
  - instr(o) $\in$ **N** instruction that implements operation o
  - cycle(o) $\in$ **N** issue cycle for operation o
  - active(o) $\in$ {0,1} whether operation o is active

- Derived variables
  - start(t) start of live range of temporary t

    = cycle(o) where o defines t
  - end(t) end of live range of temporary t

    = max { cycle(o) | o uses t }

Sanity Constraints

- Copy operation *o* is active ⇔ no coalescing

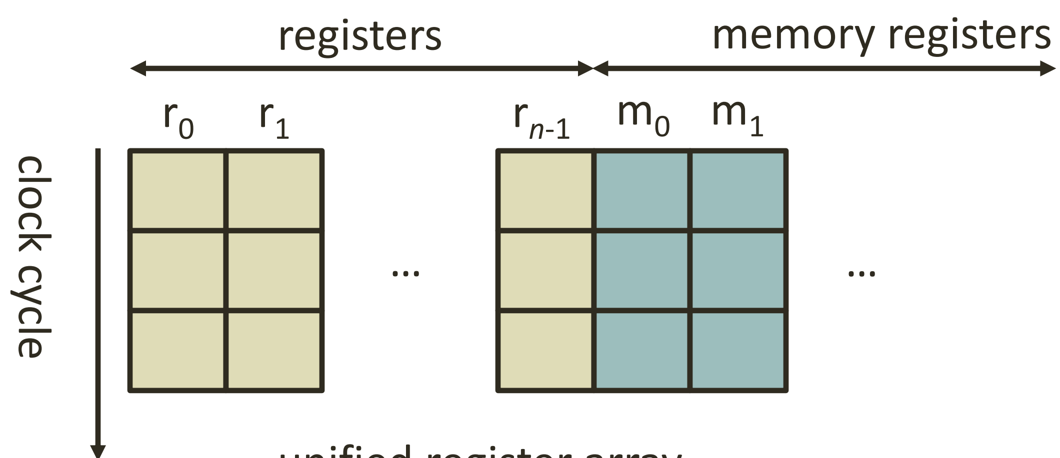
  $\text{active}(o) = 1 \Leftrightarrow \text{reg}(s) \neq \text{reg}(d)$

  - *s* is source of move, *d* is destination of move operation *o*

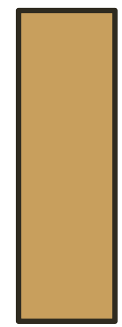
- Operations implemented by suitable instructions
  - single possible instruction for non-copy operations

- Miscellaneous
  - some registers are pre-assigned
  - some instructions can only address certain registers (or memory)

Geometrical Interpretation



unified register array

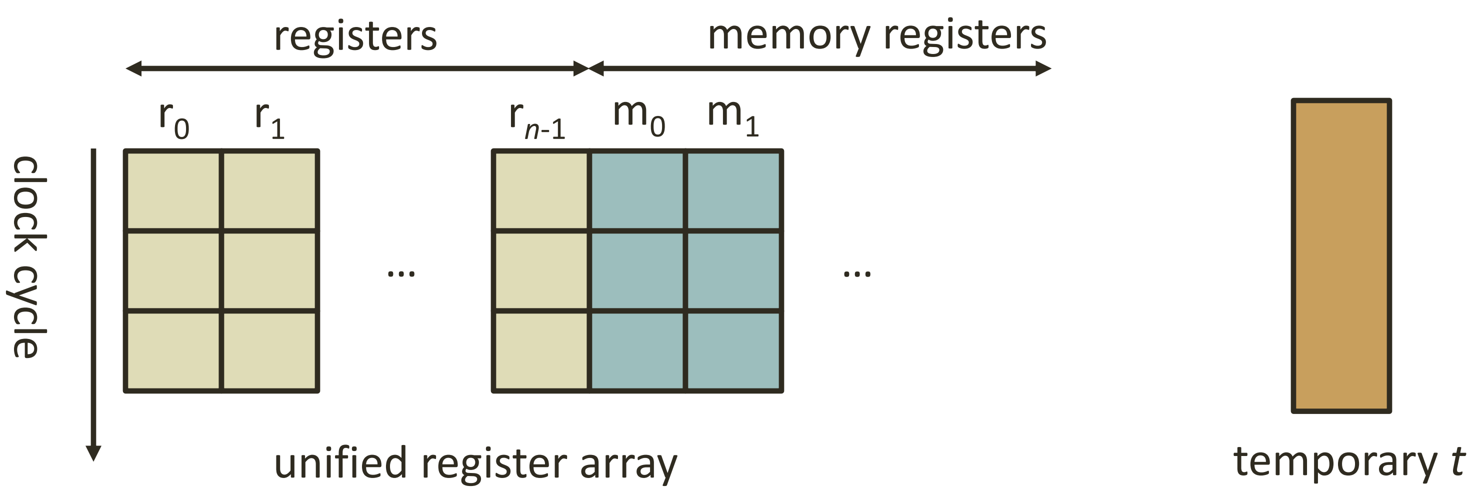


temporary *t*

- Temporary *t* is rectangle
  - width is 1 (occupies one register)
  - top = start(*t*) issue cycle of def
  - bottom = end(*t*) last issue cycle of any use

- Consequence of linear live range (basic block + SSA)

Register Assignment



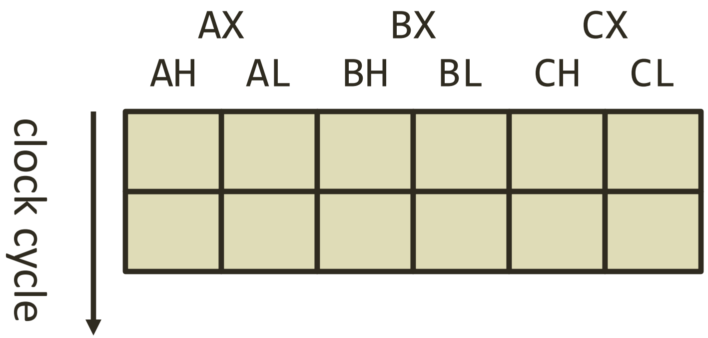
- Register assignment = geometric packing problem
  - find horizontal coordinates for all temporaries
  - such that no two rectangles for temporaries overlap
- For block *B*

  $$\text{nooverlap}(\{\langle \text{reg}(t), \text{reg}(t)+1, \text{start}(t), \text{end}(t)\rangle \mid t \in B\})$$

Register Packing

- Temporaries might have different width width(t)
  - many processors support access to register parts
  - still modeled as geometrical packing problem [Pereira & Palsberg, 2008]

Register Packing



width(t_1)=1

width(t_3)=2

width(t_3)=1

width(t_4)=2

- Temporaries might have different width   width(t)
  - many processors support access to register parts
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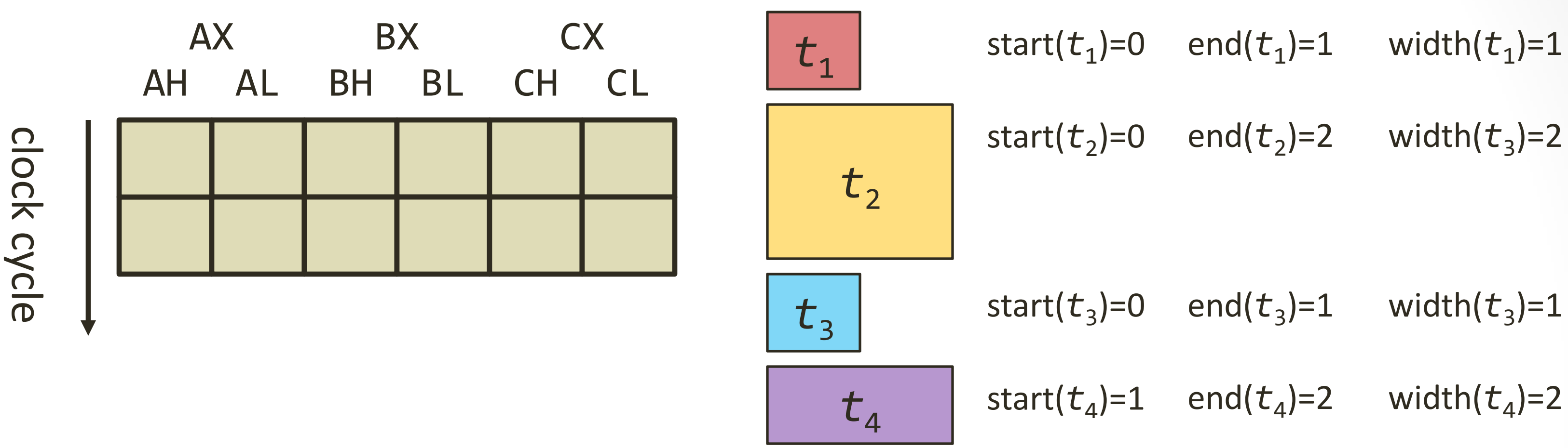
- Example: Intel x86
  - assign two 8 bit temporaries (width = 1) to 16 bit register (width = 2)
  - register parts: AH, AL, BH, BL, CH, CL
  - possible for 8 bit: AH, AL, BH, BL, CH, CL
  - possible for 16 bit: AH, BH, CH

Register Packing



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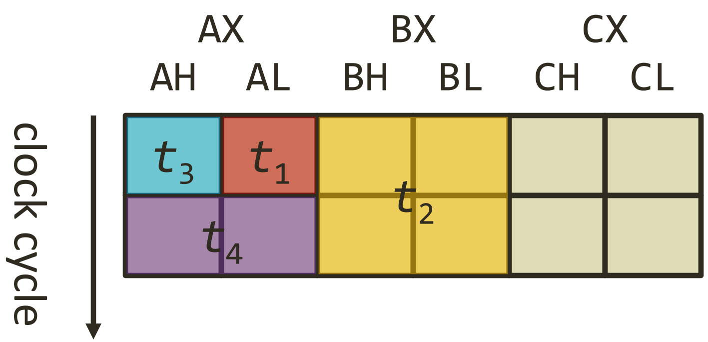

start(t_1)=0 end(t_1)=1 width(t_1)=1

start(t_2)=0 end(t_2)=2 width(t_3)=2

start(t_3)=0 end(t_3)=1 width(t_3)=1

start(t_4)=1 end(t_4)=2 width(t_4)=2

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  - many processors support access to register parts
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  - possible for 16 bit: AH, BH, CH

Modeling Register Packing

- Take width of temporaries into account (for block B)

  $\text{nooverlap}(\{\langle \text{reg}(t),\text{reg}(t)+\text{width}(t),\text{start}(t),\text{end}(t)\rangle \mid t \in B\})$

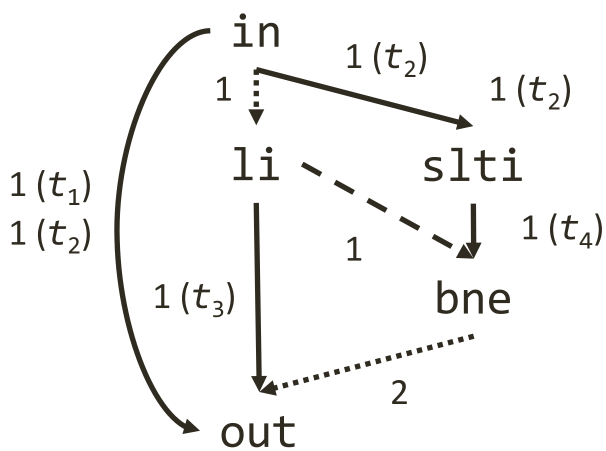
- Exclude sub-registers depending on width(t)
  - simple domain constraint on reg(t)

Local instruction scheduling (standard)

INSTRUCTION SCHEDULING

Dependencies

t_3←li
t_4←slti t_2
bne t_4



- Data and control dependencies
  - data, control, artificial (for making in and out first/last)

- If operation o_2 depends on o_1:

$$\text{active}(o_1) \land \text{active}(o_2) \rightarrow$$
$$\text{cycle}(o_2) \geq \text{cycle}(o_1) + \text{latency}(\text{instr}(o_1))$$

Processor Resources

- Processor resources: functional units, data buses, ...
  - also: instruction bundle width for VLIW processors (how many instructions can be issued simultaneously)

- Classical cumulative scheduling problem — functional units
  - processor resource has capacity — #units
  - instructions occupy parts of resource — 1 unit
  - resource consumption can never exceed capacity

- Modeling for block *B*

  $$\text{cumulative}(\{\langle \text{cycle}(o), \text{dur}(o,r), \text{active}(o) \times \text{use}(o,r) \rangle \mid o \in B\})$$

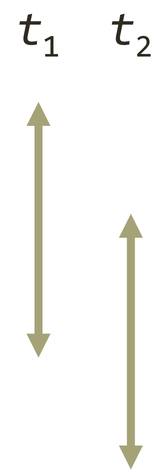
Ultimate Coalescing & Spill Code Optimization

using alternative temporaries

ADVANCED REGISTER ALLOCATION

Interference Too Naïve!

```
t1 ←     …
    …
t2 ← mv t1
    …
…  ← … t1 …
    …
…  ← … t2 …
```



t_1 and t_2 interfere

- Move-related temporaries might interfere…

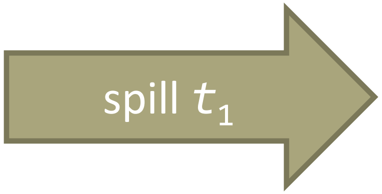
  …but contain the same value!
- Ultimate notion of interference =

  temporaries interfere ⇔ their live ranges overlap and **they have different values**

[Chaitin ea, 1981]

Spilling Too Naïve!

$t_1 \leftarrow \quad \ldots$
$\quad\ldots$
$\ldots \leftarrow \ldots\, t_1 \ldots$
$\quad\ldots$
$\ldots \leftarrow \ldots\, t_1 \ldots$



$t_1 \leftarrow \quad \ldots$
$t_2 \leftarrow \texttt{st}\ t_1$
$\quad\ldots$
$t_3 \leftarrow \texttt{ld}\ t_2$
$\ldots \leftarrow \ldots\, t_3 \ldots$
$\quad\ldots$
$t_4 \leftarrow \texttt{ld}\ t_2$
$\ldots \leftarrow \ldots\, t_4 \ldots$

- Known as **spill-everywhere** model
  - reload from memory before every use of original temporary
- Example: t_3 should be used rather than reloading t_2
  - t_2 allocated in memory!

Alternative Temporaries

- Used to track which temporaries are equal
- Representation is augmented by operands
  - act as def and use ports in operations
  - temporaries hold values transferred among operations by connecting to operands
- Example
  - operation $t_2 \leftarrow \texttt{abs}\ t_1$
  - transformed to $p_2{:}t_2 \leftarrow \texttt{abs}\ p_1{:}t_1$ (p_1, p_2 operands)
  - if t_1 and t_3 hold same value then transformed to

    $p_2{:}t_2 \leftarrow \texttt{abs}\ p_1{:}\{t_1,t_3\}$

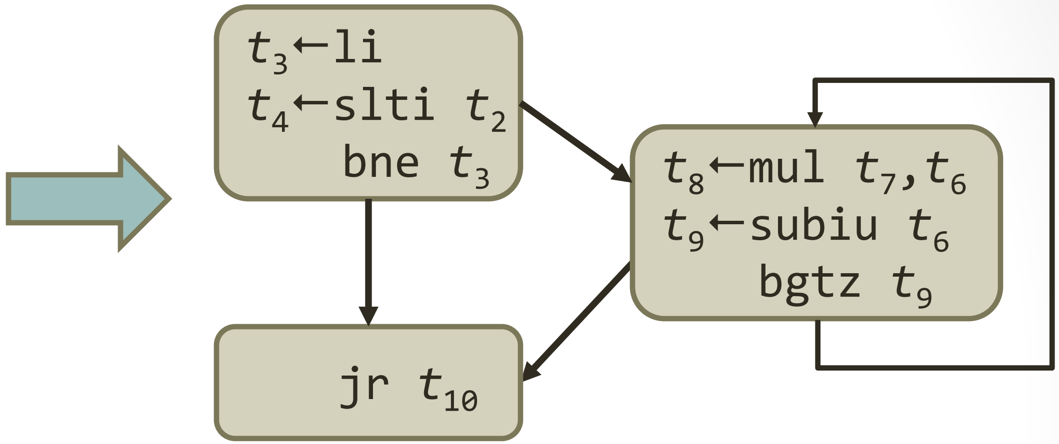
    where either t_1 or t_3 can be connected to p_1
- Model: whether a temporary is live (it is being used)

Register allocation for entire functions

GLOBAL REGISTER ALLOCATION

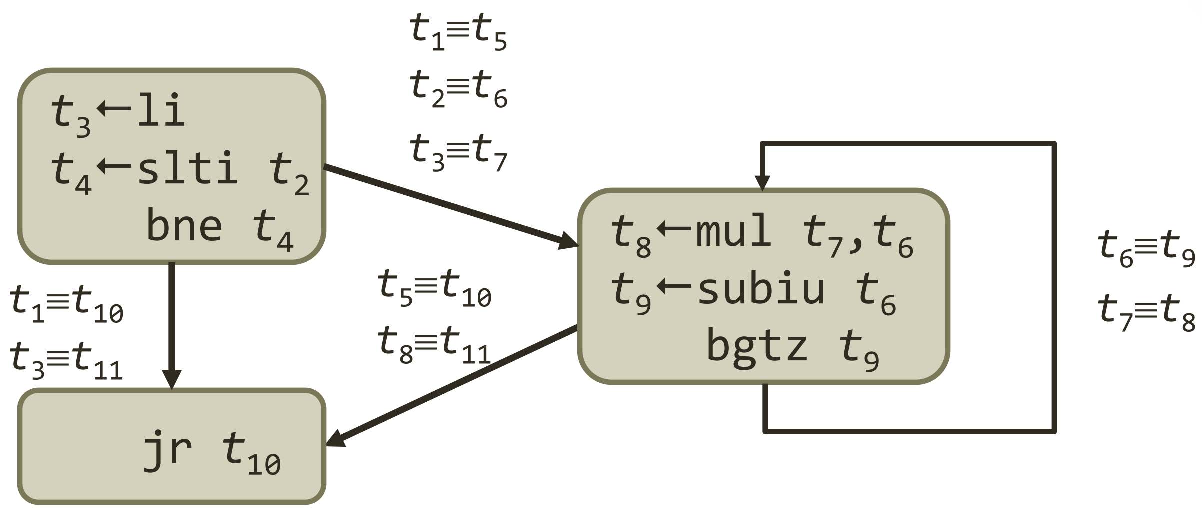
Entire Functions

```
int fac(int n) {
  int f = 1;
  while (n > 0) {
    f = f * n; n--;
  }
  return f;
}
```



- Use control flow graph (CFG) and turn it into LSSA form
  - edges = control flow
  - nodes = basic blocks (no control flow)
- LSSA = linear SSA = SSA for basic blocks plus… to be explained

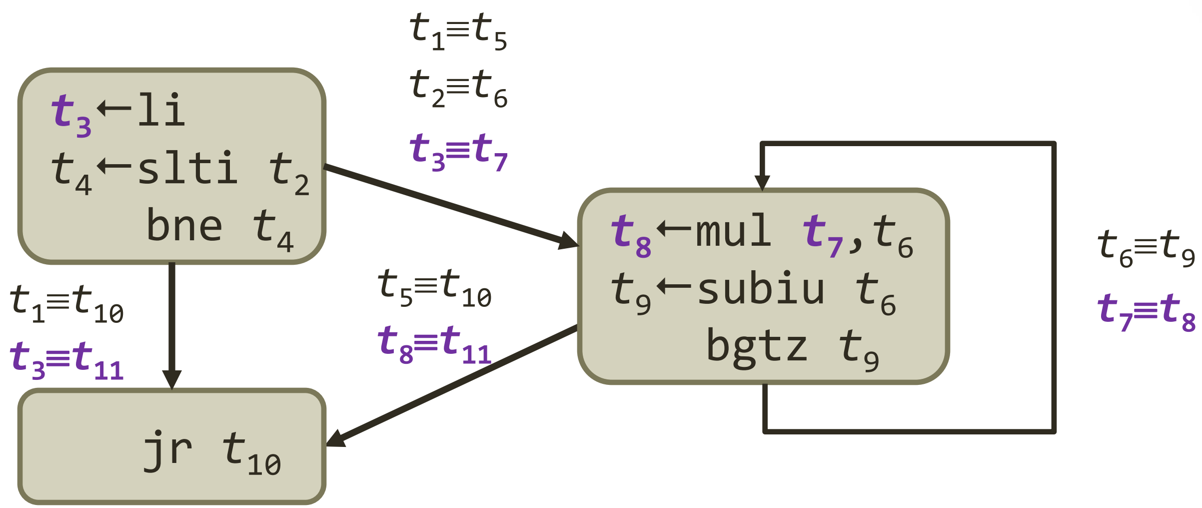
Linear SSA (LSSA)



- Linear live range of a temporary cannot span block boundaries
- Liveness across blocks defined by temporary congruence $\equiv$

$t \equiv t' \quad \Leftrightarrow \quad$ represent same original temporary

Linear SSA (LSSA)

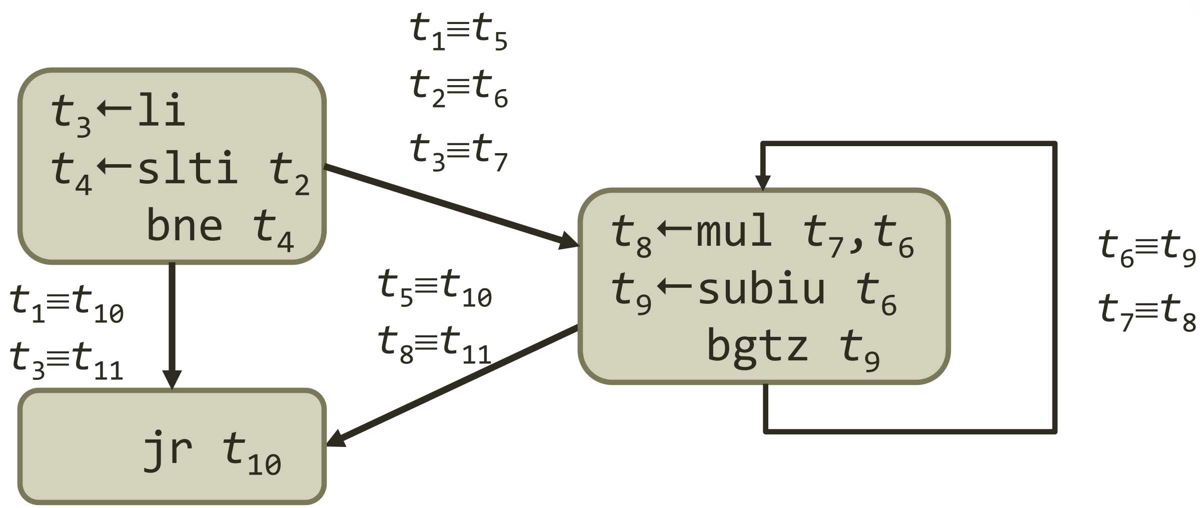


- Linear live range of a temporary cannot span block boundaries
- Liveness across blocks defined by temporary congruence ≡

  $t \equiv t'$ ⇔ represent same original temporary

- Example: t_3, t_7, t_8, t_{11} are congruent
  - correspond to the program variable f (factorial result)
  - not discussed: t_1 return address, t_2 first argument, t_{11} return value

Linear SSA (LSSA)



- Linear live range of a temporary cannot span block boundaries
- Liveness across blocks defined by temporary congruence ≡

  $t \equiv t'$ ⇔ represent same original temporary
- Advantage
  - simple modeling for linear live ranges (geometrical interpretation)
  - enables problem decomposition for solving

Global Register Allocation

- Try to coalesce congruent temporaries
  - this is why coalescing is (even more) crucial in this model
- Introduces natural problem decomposition
  - master problem (function)   coalesce congruent temporaries
  - slave problems (basic blocks) register allocation & instruction scheduling
- What is happening
  - if register pressure is low…

    no copy instruction needed (nop)

    = coalescing
  - if register pressure is high…

    copy operation might be implemented by a move

    = no coalescing

    copy operation might be implemented by a load/store

    = spill

DISCUSSION

Solving

- Approach
  - use master-slave decomposition
  - use naïve (very) portfolio of heuristics for basic blocks
  - use some pre-solving (symmetry, no-goods, dominance)
  - not very advanced (future work)

- Benchmark setup
  - selection of medium-sized functions (25 to 1000 instructions)
  - comparison to LLVM 3.3 for Qualcomm's Hexagon V4 using –O3
  - run for ten iterations where each iteration is given more time
  - using Gecode 4.2.1
  - full details in [Castañeda ea, LCTES 2014]

Experiments Summary

- Code quality (estimated)
  - 7% mean improvement over LLVM
  - provably optimal for 29% of functions
- Quadratic average (roughly) complexity up to 1000 instructions
- Can be easily changed to optimize for code size
  - 1% mean improvement over LLVM

Related Approaches

- Idea and motivation in Unison for combinatorial optimization is absolutely not new!
  - starting in the early 1990s
    [Castañeda & Schulte, CoRR 2014]
- Approaches differ
  - which code generation tasks covered
  - which technology used (ILP, CP, SAT, Genetic Algorithms, ...)
- Common to most approaches
  - compilation unit is basic block, or
  - just a single task covered, or
  - very poor scalability
- Challenge: integration, robustness, and scalability

Unique to Unison Approach

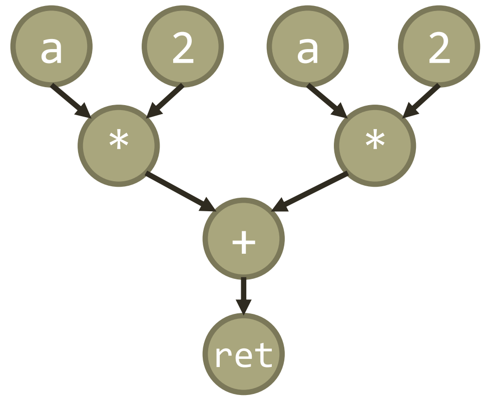
- First global approach for register allocation (function as compilation unit)
- Constraint programming using global constraints
  - sweet spot: cumulative and nooverlap
- Full register allocation with ultimate coalescing, packing, spilling, and spill code optimization
  - key property of model: spilling is internalized
- Robust at the expense of optimality
  - problem decomposition
- But: instruction selection not yet there!

Instruction Selection

[Based on slides from Gabriel Hjort Blindell]

Graph-based Instruction Selection

```
int f(int a) {
  int b = a * 2;
  int c = a * 4;
  return b + c;
}
```



- Represent program as graph **program graph**

Graph-based Instruction Selection

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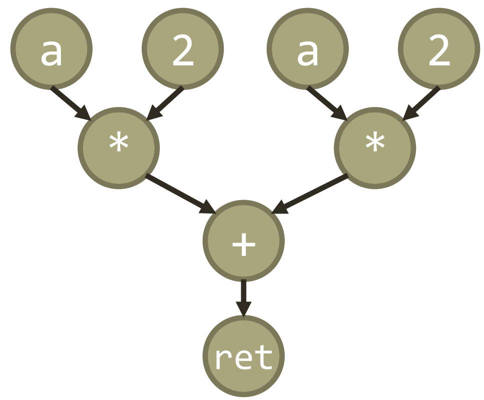

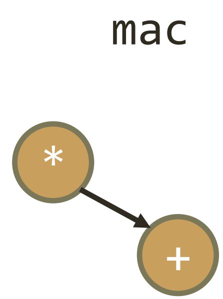

- Represent program as graph **program graph**
- Represent instructions as graph **instruction graph**

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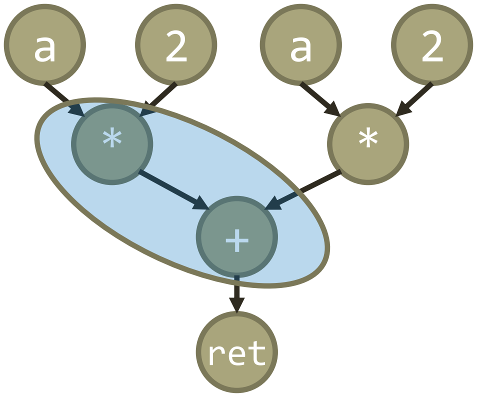
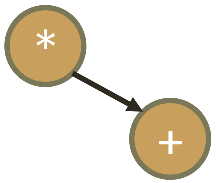


mac



- Represent program as graph **program graph**
- Represent instructions as graph **instruction graph**
- Select **matches** such that program graph is **covered**

Graph-based Instruction Selection

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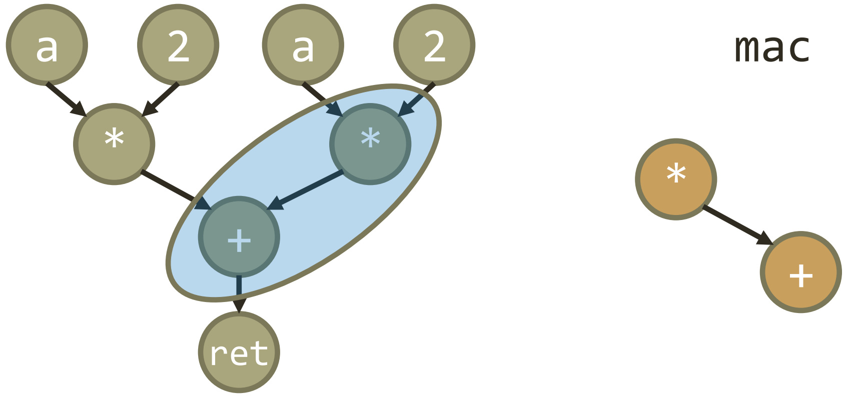


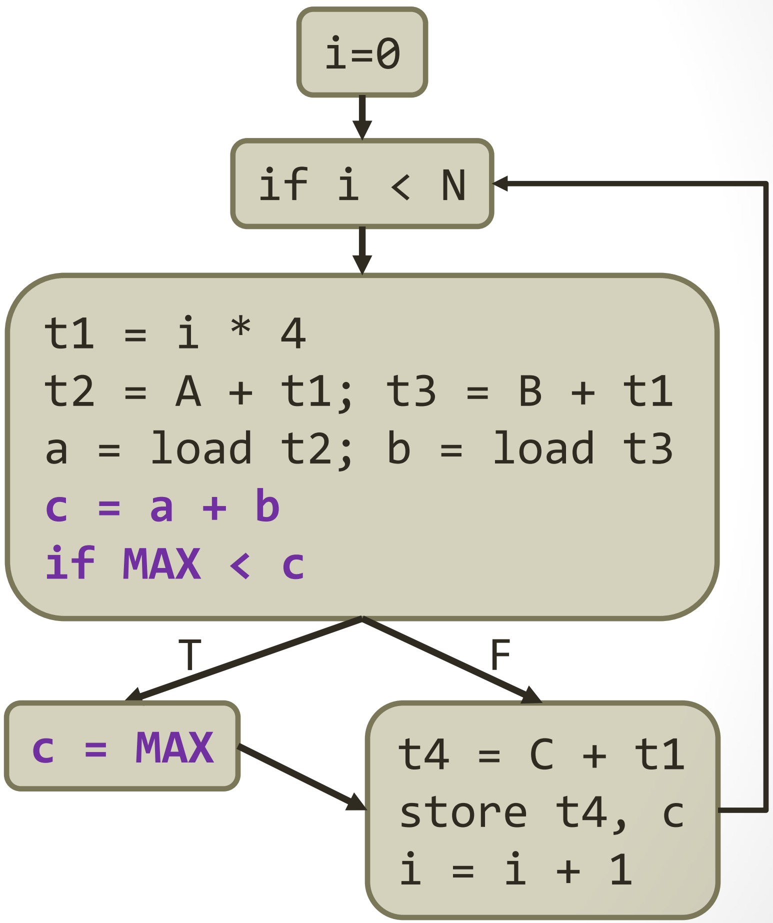
- Represent program as graph **program graph**
- Represent instructions as graph **instruction graph**
- Select **matches** such that program graph is **covered**

State of the Art

- Local instruction selection
- Program graphs per block
- Graphs restricted to data flow
  - cannot handle control flow such as branching instructions
- Greedy heuristics
  - For example, maximal munch

Instruction Examples

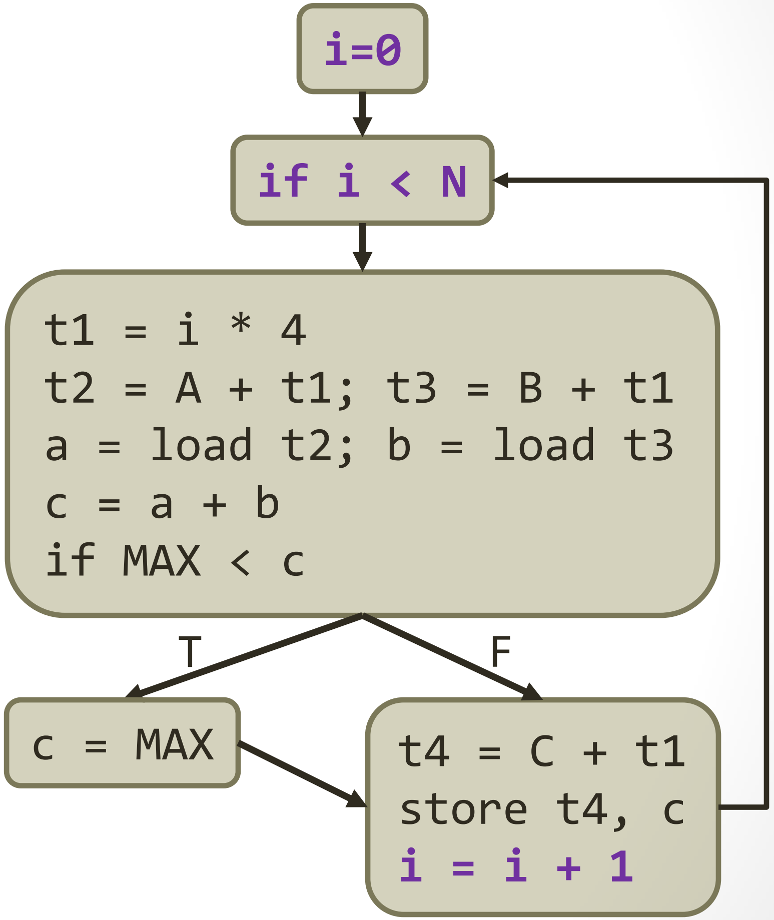
- **satadd**



- Exists in many DSPs
- Incorporates control flow
- Extends across basic blocks

Instruction Examples

- `satadd`
- **`repeat`**

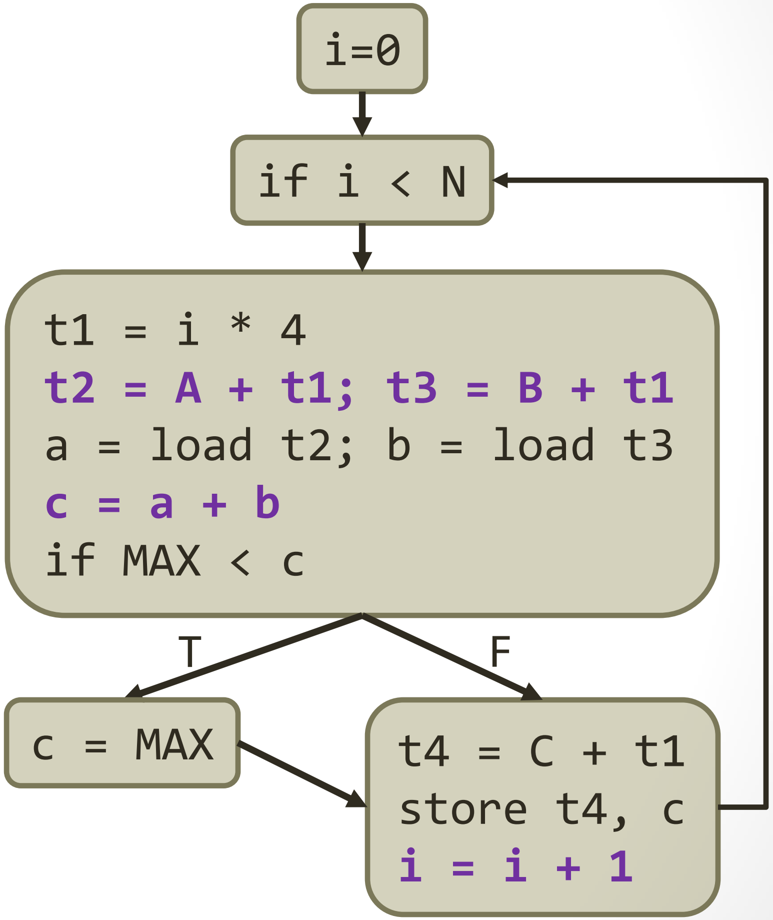


- Exists in many processors
  - for example Intel’s x86
- Incorporates control flow
- Extends across basic blocks

Instruction Examples

- satadd
- repeat
- **add4**

- SIMD-style instruction
  - very common
- Requires **global code motion**
  - move computations across blocks
- Depending on hardware may require copying
  - different register file

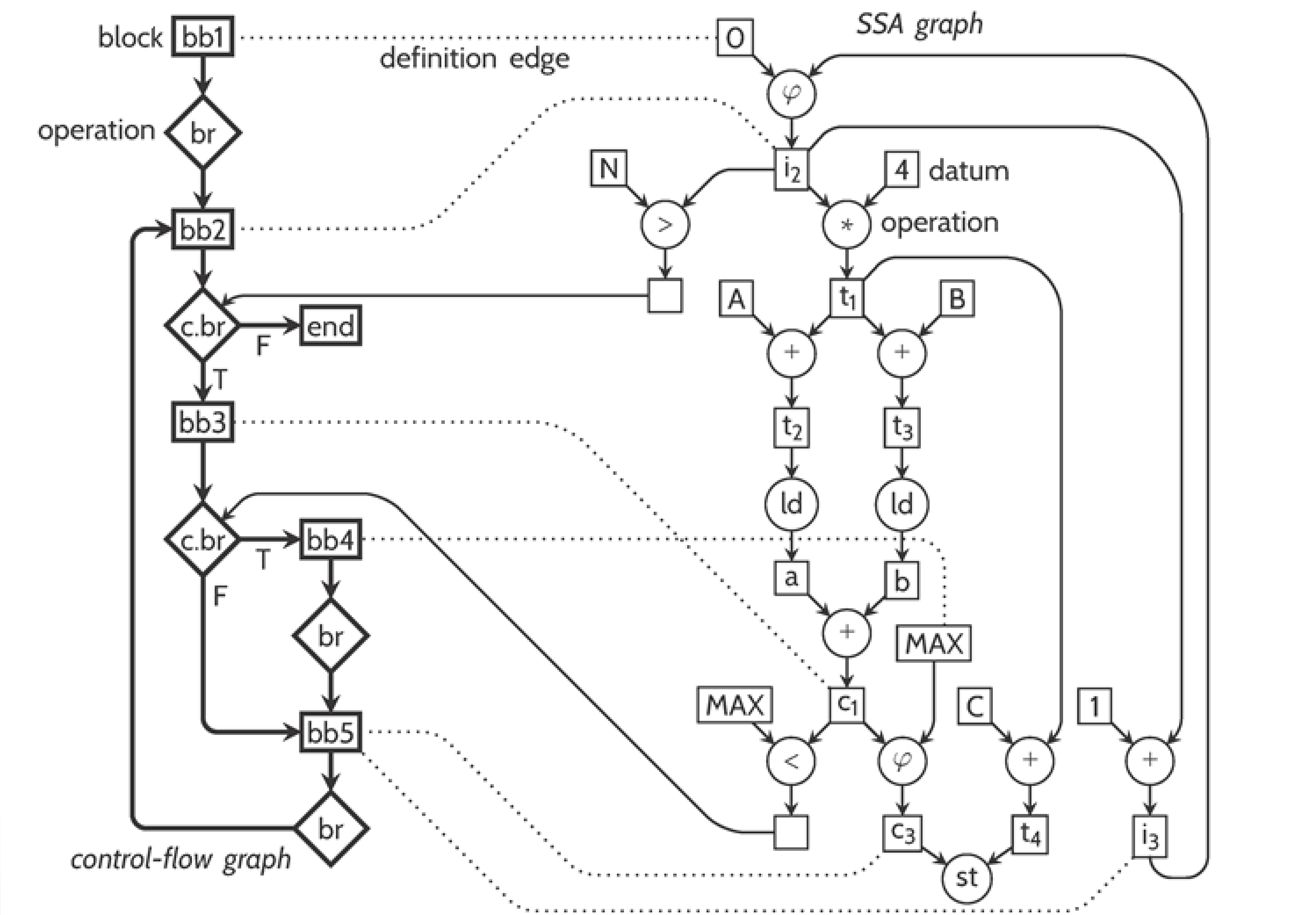


Universal Instruction Selection

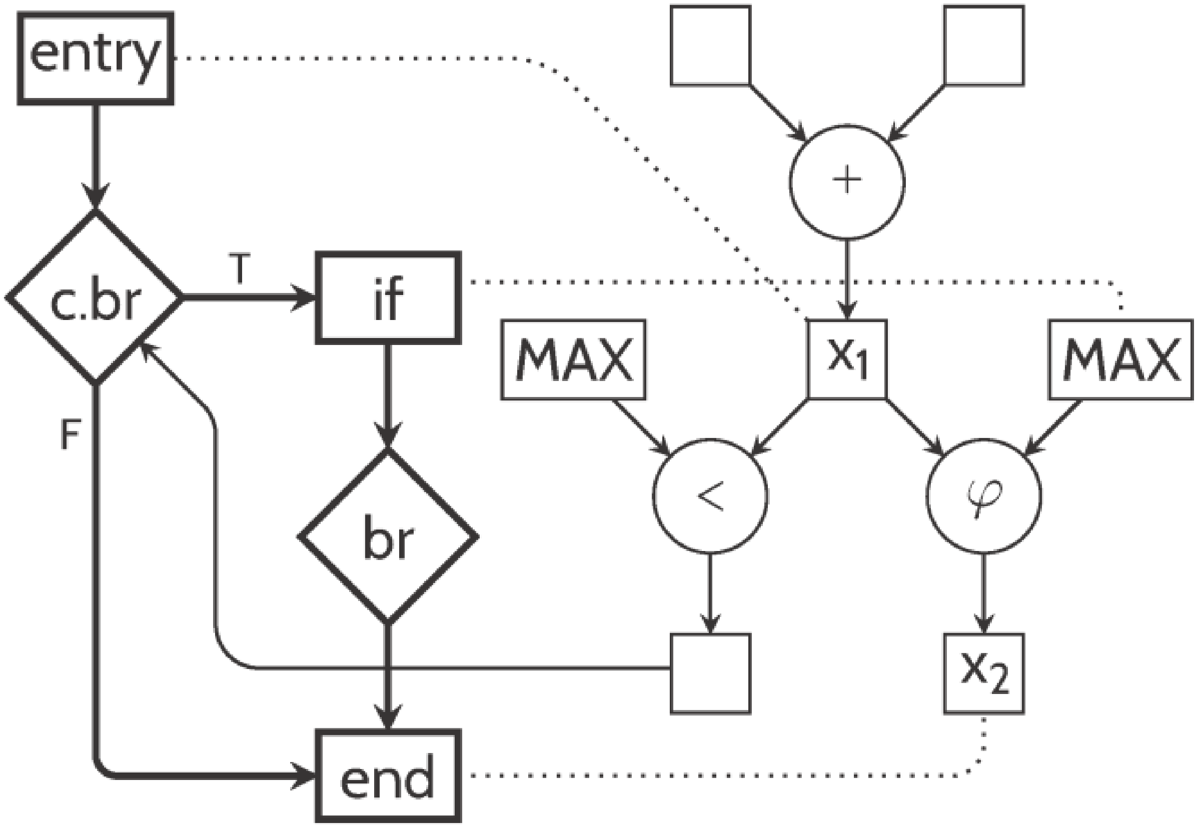
- Global instruction selection
- Program graphs for entire functions
- Instruction graphs capture both data and control flow
    - handles broad range of instructions found in today's processors
- Integrates global code motion
- Takes data-copying overhead into account

- Presupposes an expressive approach such as CP

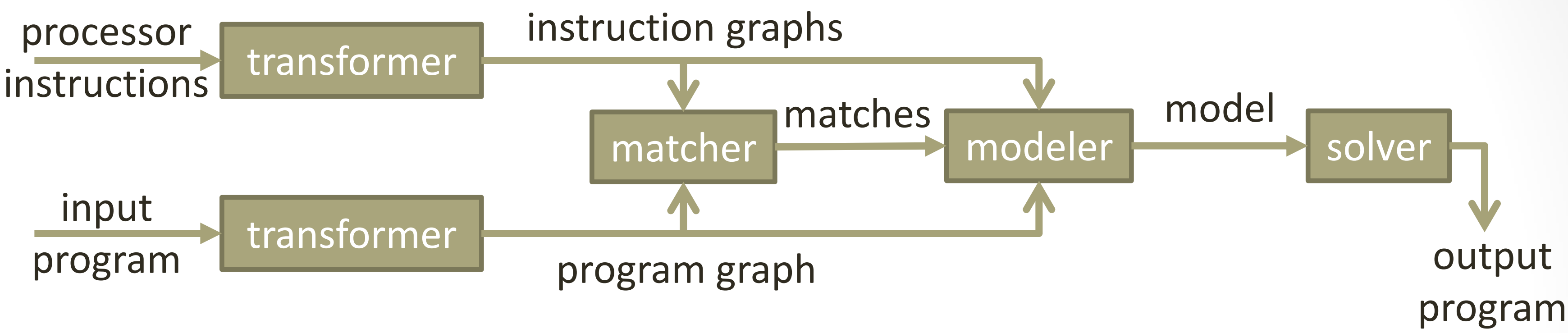
Program Graph (Example)



Instruction Graph (`satadd`)



Approach



- Before: create instruction graphs
- Code generation
  - create program graph
  - compute possible matches (standard algorithm VF2 [Cordella ea, 2004])
  - generate model in MiniZinc
  - solve model with CPX 1.0.2

Model Summary

- Decision variables
  - which match is selected?
  - in which block are selected matches placed?
  - in which block is data made available?
- Constraints (selection)
  - operations must be covered by exactly one match
  - control flow cannot be moved
  - data must be defined before used
  - definition edges must be enforced
  - blocks must be ordered (respect fall-through branching if possible)
  - implied and dominance constraints
- Objective functions
  - minimize estimated execution time
  - minimize code size

Experiments

- Benchmarks
  - 16 functions from MediaBench
  - program graphs have 34-203 nodes
  - all models solved to optimality with CPX 1.0.2
- For Simple MIPS32
  - simple RISC architecture: worst-case scenario
  - surprise: 1.4% mean speedup over LLVM 3.4
  - better: global code motion; worse: constant reloading
  - runtimes: 0.3-83.2 seconds, median 10.5 seconds
- For Funky MIPS32 (made up)
  - MIPS32 + common SIMD instructions: good case
  - 3% mean speedup over Simple MIPS32
  - surprise: sometimes SIMD-style is not really that good!
  - runtimes: 0.3-146.8% seconds, median 10.5 seconds

Discussion

- Overcomes many restrictions of state-of-the-art approaches
  - control flow
  - global code motion
  - sophisticated instructions

- Model and representation designed together
  - expressive representation requires expressive models

- Limitations
  - constant reloading
  - if-conversion (predication), well: no approach can do this anyway!

SUMMARY

Now and Then…

- Status
  - instruction scheduling: local, standard
  - register allocation: global, unique
  - instruction selection: global, unique
  - not fully integrated
  - solving pretty naïve
- Future
  - instruction scheduling: superblocks, if-conversion (predication)
  - register allocation: rematerialization (done)
  - more sophisticated solving
  - integration!!!

Project & Goals

- Unison has a considerable engineering part
  - processor descriptions (separate large project)
  - robust and maintainable tool chain
  - testing and transfer
- A production-quality tool that will be deployed
  - industrial strength re-implementation started
- An open-source contribution to LLVM
  - repositories on GitHub since October 2016
  - need to convince LLVM developers…
- Real significance

simplicity even for today's freak processors